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                                                         F. Michel, Ed.
                                                             UCLouvain
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Coding and congestion control in transport
draft-kuhn-coding-congestion-transport-00

Abstract

   There are discussions on how loss-based congestion controls consider
   lost packets that have been recovered by a coding mechanism. This
   document analyses to what extent transport protocols could ignore
   such signals and proposes best current practices on the interaction
   between congestion control and coding mechanism at the transport
   layer. Coding for tunnels is out-of-the scope of the document.
   Examples of interest for the proposed solution is to better deal with
   tail losses or with networks with non-congestion losses.

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Table of Contents

1. Introduction

[RFC5681] defines TCP as a loss-based congestion control. Coding
mechanisms can be deployed and could hide congestion signals to the
sender.

Because loss-based and delay-based congestion controls are deployed
in the current Internet, this memo discusses simple best practices on
how coding and congestion control mechanisms could coexist.

The proposed recommendations apply for coding at the transport layer
(coding for tunnels is out-of-the scope of the document). Examples
of interest for the proposed solution is to better deal with tail
losses or with networks with non-congestion losses.

2. Base solution description

The base solution can be described as follows:

o The receiver MUST indicate to the sender that one or multiple
   packets have been recovered using a coding scheme. Such "repaired
   packet signal" could be based on existing signals (even if the
   existing signal was not designed for that purpose, such as ECN) or
   on new type of signals (such as a RECOVERED frame in QUIC).

o The sender MUST be able to detect the "recovered packet signal".
   The base solution does not describe how the sender reacts to such
   signal.

The proposed solution applies for coding in the transport layer. The
proposed approach is inline with the one in
[I-D.swett-nwcrg-coding-for-quic].

The proposed solution does not applies for the interaction between
coding under the transport layer (i.e. not end-to-end), such as
coding for tunnels.

3. Sender-side coding solutions

This section presents solutions for a sender to add application
coding.

3.1. Coded packets without considering CWND progression

In this solution, the coded packets are sent on top of what is
allowed by a congestion window. Examples of the solution could be
adding a given pourcentage of the congestion window as supplementary
packets or sending a given amount of coded packets at a given rate.
The redundancy flow can be decorrelated from the congestion control
that manages source packets : a secondary congestion control can be
introduced, such as in coupled congestion control for RTP media
[I-D.ietf-rmcat-coupled-cc]. An example would be to exploit a lower
than best-effort congestion control [RFC6297].

The advantage of such solution is that coding would help in
challenges cases where transmission losses are persistent.

The drawback of such solution is that it may result in coding
solutions being unfair towards non-coding solutions. This solutions
may result in adding congestion in congested networks.

3.2. Coded packets driven by CWND progression

In this solution, the coded packets are sent within what the
congestion window allows, such as in [CTCP]. Examples of the
solution would be sending coded packets when there is no more data to

transmit or preferably send coded packets instead of the following packets in the send buffer.

The advantage of this solution is that it does not contribute in adding more congestion than the congestion window allows. Indeed, all traffic (source and redundancy) is controlled by one congestion control only and TCP metrics for fairness can be indifferently applied in this case.

The main drawback is the decrease of goodput if coded packets are sent but are not used at the client side.

4. Sender-side reaction to recovered packet signals

Delay-based congestion controls ignore packets that have been repaired with coding. There is no need to define best current pratices in this case. However, more discussions are required for congestion controls that use loss as congestion signals (potentially among other congestion detection mechanism).

4.1. The sender congestion control considers recovered packet signals as congestion-implied packet losses

In this solution, the sender reacts to recovered packet signals as to congestion-implied packet losses. That being said, this does not necessarily means that the packets have actually been lost. The server may have other means to identify that the packet was just out-of-ordered and ignore the recovered packet signals.

The advantages of the solution are (1) that coding mechanisms do not hide congestion signals, such as packets voluntary dropped by a AQM [RFC7567] and (2) packets may be recovered faster than with traditionnal retransmission mechanisms.

The drawback of this solution is that, if there is a high non-congestion loss rate, the congestion control throughput may decrease drastically. Reporting this amount of loss to a sender may reduce the application goodput when there is no actual congestion.

4.2. The sender adapts its window reduction to recovered packet signals

In this solution, the sender does not reduce the congestion window with the same amount when the "recovered packet signal" is received, i.e. when a packet has been lost but recovered. Example of this solution could be based on [RFC8511] or considering that recovering an isolated packet is not an actual sign of congestion.

The advantage of the solution is that in cases where there is no actual congestion, coding could help in improving the transmission without ignoring congestion signals.

The main drawback is the precised design of the solution and its interaction with AQM mechanisms [RFC7567]. Moreover there may be fairness issues since AIMD convergence may not be guaranteed.

4.3. The sender ignores recovered packet signals

This is the case for delay-based congestion controls. The interaction between delay-based congestion controls and the delay induced by a coding mechanisms is an open research activity. That being said, a potential approach would be that loss-based congestion control ignores the "recovered packet signal".

The advantage of this solution is that coding would provided substantial benefits in cases where there are transmission losses.

However, this solution hides potential congestion losses, making it unfair to other congestion controls.

5. Summary

This section provides a summary on the content in previous sections. The Figure 1 sums up some recommendations. It is worth pointing out that the "coding without congestion" considers that coded packets are sent along with original data packets, in opposition with the solution where coded packets are transmitted only when there is no more original packets to transmit. Moreover, the values indicated in this Figure consider a channel that does not exhibit a high loss pattern.

Sender-side reaction to recovered packet signals	Sender-side coding solutions	
	Coding adding congestion	Coding without congestion
React as loss	fairness: ~ real-time: + bulk: ~	fairness: ++ real-time: + bulk: -
Adapt window reduction	fairness: ~ real-time: + bulk: +	fairness: + real-time: + bulk: -
Ignore signals	fairness: - real-time: + bulk: +	fairness: - real-time: + bulk: -

Figure 1: Recommendations

6. Acknowledgements

Many thanks to TBD.

7. Contributors

TBD

8. IANA Considerations

This memo includes no request to IANA.

9. Security Considerations

Security section goes here.

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Authors' Addresses

   Nicolas Kuhn (editor)
   CNES

   Email: nicolas.kuhn@cnes.fr

   Emmanuel Lochin (editor)
   ISAE-SUPAERO

   Email: emmanuel.lochin@isae-supaero.fr

   Francois Michel (editor)
   UCLouvain

   Email: francois.michel@uclouvain.be