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Use case for a scalable and topology aware MPLS data plane monitoring system draft-geib-spring-oam-usecase-02

Abstract

This document describes features and a use case of a path monitoring system. Segment based routing enables a scalable and simple method to monitor data plane liveliness of the complete set of paths belonging to a single domain. Compared with legacy MPLS ping and path trace, MPLS topology awareness reduces management and control plane involvement of OAM measurements while enabling new OAM features.

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1. Introduction

It is essential for a network operator to monitor all the forwarding paths observed by the transported user packets. The monitoring flow must be forwarded in dataplane in a similar way as user packets. Problem localization is required.

This document describes a solution to this problem statement and illustrates it with use-cases.

The solution is described for a single IGP MPLS domain.

The solution applies to monitoring of LDP LSP's as well as to monitoring of Segment Routed LSP's. Segment Routing simplifies the solution by the use of IGP-based signalled segments as specified by [ID.sr-isis]. Thus a centralised monitoring unit is MPLS topology aware in a Segment Routed domain and this topology awareness is used for OAM purposes. The MPLS path monitoring system described by this document can be realised with pre-Segment based Routing (SR) technology. Making such a monitoring system aware of a domains complete MPLS topology requires e.g. management plane access. To avoid the use of stale MPLS label information, IGP must be monitored and MPLS topology must be timely aligned with IGP topology. Obviously, enhancing IGPs to exchange of MPLS topology information significantly simplifies and stabilises such an MPLS path monitoring system.

This document adopts the terminology and framework described in [ID.sr-archi]. It further adopts the editorial simplification explained in section 1.2 of the segment routing use-cases [ID.sr-use].

The proposed solution offers several benefits for network monitoring. A single centralized monitoring device is able to monitor the complete set of a domains forwarding paths. OAM packets never leave data plane. Legacy path trace is still required. In addition to Segment Routing related IGP extensions, also RFC 4379 features should be extended to support detection of SR routed paths. They further should be enhanced to support all deployed IP/MPLS entropy options. In an IPv6 domain, a MPLS like tree trace functionality is desirable.

Faults can be localized:

- o by IGP LSA analysis.
- o by correlation between different probes.

o by MPLS traceroute and adapted ping messages.

The proposed solution requires topology awareness as well as a suitable security architecture. Topology awareness is an essential part of link state IGPs. Adding MPLS topology awareness to an IGP speaking device hence enables a simple and scaleable data plane monitoring mechanism.

MPLS OAM offers flexible features to recognise an execute data paths of an MPLS domain. By utilsing the ECMP related tool set of RFC 4379 [RFC4379], a segment based routing LSP monitoring system may:

- o easily detect ECMP functionality and properties of paths at data level.
- o construct monitoring packets executing desired paths also if ECMP is present.
- o limit the MPLS label stack of an OAM packet to a minmum of 3 labels.

MPLS OAM supports detection and execution of ECMP paths quite smart. This document is foscused on MPLS path monitoring.

Alternatively, any path may be executed by building suitable label stacks. This allows path execution without ECMP awareness.

The MPLS path monitoring system may be a specialised system residing at a single interface of the domain to be monitored. As long as measurement packets return to this or another interface to a specialised OAM system, the MPLS monitoring system is the single entity pushing monitoring packet label stacks. Concerns about router label stack pushing capabilities don't apply in this case.

First drafts discussing requirements, extensions of $\frac{RFC4379}{AFC4379}$ and possible solutions to allow SR usage as described by this document are at hand, see $\frac{[ID.sr-4379ext]}{AFC4379ext}$ and $\frac{[ID.sr-0am_detect]}{AFC4379ext}$.

2. An MPLS topology aware path monitoring system

An MPLS path monitoring system (PMS) which is able to learn the IGP LSDB (including the SID's) is able to build a measurement packet which executes every arbitrary chain of paths. A node connected to an SR domain is MPLS topology aware (the node knows all related IP adresses, MPLS SIDs and labels).

Let us describe how the PMS can check the liveliness of the MPLS

transport path between LER i and LER j and then monitor it.

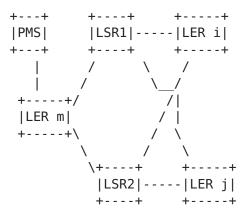
The PMS may do so by sending packets carrying the following MPLS label stack infomation:

- o Top Label: a path from PMS to LER i This is expressed as Node SID of LER i.
- o Next Label: the path that needs to be monitored from LER i to LER j. If this path is a single physical interface (or a bundle of connected interfaces), it can be expressed by the related AdjSID. If the shortest path from LER i to LER j is supposed to be monitored, the Node-SID (LER j) can be used. Another option is to insert a list of segments expressing the desired path (hop by hop as an extreme case). If LER i pushes a stack of Labels based on a SR policy decision and this stack of LSPs is to be monitored, the PMS needs an interface to collect the information enabling it to address this SR created path.
- o Next Label or address: the path back to the PMS. Likely, no further segment/label is required here. Indeed, once the packet reaches LER j, the 'steering' part of the solution is done and the probe just needs to return to the PMS. This is best achieved by popping the MPLS stack and revealing a probe packet with PMS as destination address (note that in this case, the source and destination addresses could be the same). If an IP address is applied, no SID/label has to be assigned to the PMS (if it is a host/server residing in an IP subnet outside the MPLS domain).

Note: if the PMS is an IP host not connected to the MPLS domain, the PMS can send its probe with the list of SIDs/Labels onto a suitable tunnel provding an MPLS access to a router which is part of the monitored MPLS domain.

3. SR based OAM use case illustration

3.1. Use-case 1 - LSP dataplane liveliness detection and monitoring



Example of a PMS based LSP dataplane liveness detection and monitoring

Figure 1

For the sake of simplicity, let's assume that all the nodes are configured with the same SRGB [ID.sr-archi]. as described by section 1.2 of [ID.sr-use].

Let's assign the following Node SIDs to the nodes of the figure: PMS = 10, LER i = 20, LER j = 30.

The aim is to check liveliness of the path LER i to LER j and to monitor availability of that path afterwards. The PMS does this by creating a measurement packet with the following label stack (top to bottom): 20 - 30 - 10.

LER m forwards the packet received from the PMS to LSR1. Assuming Pen-ultimate Hop Popping to be deployed, LSR1 pops the top label and forwards the packet to LER i. There the top label has a value 30 and LER i forwards it to LER j. This will be done transmitting the packet via LSR1 or LSR2. The LSR will again pop the top label. LER j will forward the packet now carrying the top label 10 to the PMS (and it will pass a LSR and LER m).

A few observations on the example given in figure 1:

o The path PMS to LER i must be available. This path must be detectable, but it is usually sufficient to apply an SPF based path.

- o If ECMP is deployed, it may be desired to measure along both possible paths, a packet may use between LER i and LER j. To do so, in a first step the PMS sends MPLS OAM packets to execute a so called tree trace between LER i and LER j and stores the IP destination addresses required to execute each detected path. This method of dealing with load balancing paths requires the smallest label stacks if long term monitoring of paths is applied after the tree trace completion.
- o The path LER j to PMS to must be be available. This path must be detectable, but it is usually sufficient to apply an SPF based path.

Once the MPLS paths (Node SIDs) and the required IP address information has been detected, the LER i to LER j can be monitored by the PMS. Monitoring doesn't require MPLS OAM functionality, it is purely based on forwarding. To ensure reliable results, the PMS should be aware of any changes in IGP or MPLS topology. Further changes in ECMP functionality at LER i will impact results. Either the PMS should be notified of such changes or they should be limited to planned maintenance. After a topology change, MPLS OAM will be useful to detect the impact of the change.

Determining a path to be executed prior to a measurement may also be done by setting up a label including all node SIDs along that path (if LER1 has Node SID 40 in the example and it should be passed between LER i and LER j, the label stack is 20 - 40 - 30 - 10). The advantage of this method is, that it does not involve MPLS OAM functionality and it is independent of ECMP functionalities. The method still is able to monitor all link combinations of all paths of an MPLS domain. If correct forwarding along the desired paths has to be checked, RFC4739 functionality should be applied also in this case.

Obviously, the PMS is able to check and monitor data plane liveliness of all LSPs in the domain. The PMS may be a router, but could also be dedicated monitoring system. If measurement system reliability is an issue, more than a single PMS may be connected to the MPLS domain.

Monitoring an MPLS domain by a PMS based on SR offers the option of monitoring complete MPLS domains with little effort and very excellent scalability. Data plane failure detection by circulating monitoring packets can be executed at any time. The PMS further executes MPLS OAM functions everywhere in the MPLS domain. It does not require access to LSR/LER management interfaces to do so. MPLS traceroutes as specified above should be executed only during off peak times (and then with limited parallel MPLS ping/trace load).

3.2. Use-case 2 - Monitoring a remote bundle

SR based probing of all the links of a remote bundle

Figure 2

R1 adresses Lx by the Adjacency SID 99x, while R2 adresses Lx by the Adjacency SID 66(x+1).

In the above figure, the PMS needs to assess the dataplane availability of all the links within a remote bundle connected to routers R1 and R2.

The monitoring system retrieves the SID/Label information from the IGP LSDB and appends the following segment list/label stack: {72, 662, 992, 664} on its IP probe (whose source and destination addresses are the address of the PMS).

MS sends the probe to its connected router. If the connected router is not SR compliant, a tunneling technique can be used to tunnel the probe and its MPLS stack to the first SR router. The MPLS/SR domain then forwards the probe to R2 (72 is the Node SID of R2). R2 forwards the probe to R1 over link L1 (Adjacency SID 662). R1 forwards the probe to R2 over link L2 (Adjacency SID 992). R2 forwards the probe to R1 over link L3 (Adjacency SID 664). R1 then forwards the IP probe to PMS as per classic IP forwarding.

3.3. Use-Case 3 - Fault localization

In the previous example, a uni-directional fault on the middle link from R1 to R2 would be localized by sending the following two probes with respective segment lists:

- o 72, 662, 992, 664
- o 72, 663, 992, 664

The first probe would fail while the second would succeed. Correlation of the measurements reveals that the only difference is

using the Adjacency SID 662 of the middle link from R1 to R2 in the non successful measurement. Assuming the second probe has been routed correctly, the fault must have been occurring in R2 which didn't forward the packet to the interface identified by its Adjacency SID 662.

4. Failure Notification from PMS to LERi

PMS on detecting any failure in the path liveliness MAY use any outof-band mechanism to signal te\he failure to LERi. This document does not not propose any specific mechanism and Operators can choose any existing or new approach.

Alternately, the Operator may log the failure in local monitoring system and take necessary action by manual intervention.

5. Applying SR to monitor LDP paths

A SR based PMS connected to a MPLS domain consisting of LER and LSR supporting SR and LDP in parrallel in all nodes may use SR paths to transmit packets to and from start and end points of LDP paths to be monitored. In the above example, the label stack top to bottom may be as follows, when sent by the PMS:

- o Top: SR based Node-SID of LER i at LER m.
- o Next: LDP label identifying the path to LER j at LER i.
- o Bottom: SR based Node-SID identifying the path to the PMS at LER j

While the mixed operation shown here still requires the PMS to be aware of the LER LDP-MPLS topology, the PMS may learn the SR MPLS topology by IGP and use this information.

6. PMS monitoring of different Segment ID types

MPLS SR topology awareness should allow the SID to monitor liveliness of most types of SIDs (this may not be recommendable if a SID identifies an inter domain interface).

To match control plane information with data plane information, RFC4379 should be enhanced to allow collection of data relevant to check all relevant types of Segment IDs.

7. Connectivity Verification using PMS

While the PMS based use cases explained in <u>Section 3</u> is sufficient to provide Continuity check between LER i and LER j, it may not help perform connectivity verification. So in some cases like data plane programming corruption, it is possible that a transit node between LER i and LER j erroneously remove the top segment ID and forward to PMS based on bottom segment ID leading to falsified path liveliness to PMS.

There are various method to perform basic connectivity verification like intermittely setting the TTL to 1 in bottom label so LER j selectively perform connectivity verification. A detailed explanation will be added in later version.

8. Extensions of related standards

<u>RFC4379</u> functions should be extended to support Flow- and Entropy Label based ECMP. Further, an <u>RFC4379</u> like functionality may be desirable for IPv6 networks.

9. IANA Considerations

This memo includes no request to IANA.

10. Security Considerations

As mentioned in the introduction, a PMS monitoring packet should never leave the domain where it originated. It therefore should never use stale MPLS or IGP routing information. Further, asigning different label ranges for different purposes may be useful. A well known global service level range may be excluded for utilisation within PMS measurement packets. These ideas shoulddn't start a discussion. They rather should point out, that such a discussion is required when SR based OAM mechanisms like a SR are standardised.

11. Acknowledgement

The authors would like to thank Nobo Akiya for his cotribution.

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