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Segment Routing Policy Architecture

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Abstract

Segment Routing (SR) allows a headend node to steer a packet flow along any path. Intermediate per-flow states are eliminated thanks to source routing. The headend node steers a flow into an SR Policy. The header of a packet steered in an SR Policy is augmented with the ordered list of segments associated with that SR Policy. This document details the concepts of SR Policy and steering into an SR Policy.

Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

Status of This Memo

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1. Introduction

Segment Routing (SR) allows a headend node to steer a packet flow along any path. Intermediate per-flow states are eliminated thanks to source routing [I-D.ietf-spring-segment-routing].

The headend node is said to steer a flow into an Segment Routing Policy (SR Policy).

The header of a packet steered in an SR Policy is augmented with the ordered list of segments associated with that SR Policy.

This document details the concepts of SR Policy and steering into an SR Policy. These apply equally to the MPLS and SRv6 instantiations of segment routing.

For reading simplicity, the illustrations are provided for the MPLS instantiations.

2. SR Policy

An SR Policy is a framework that enables instantiation of an ordered list of segments on a node for implementing a source routing policy with a specific intent for traffic steering from that node.

The Segment Routing architecture [I-D.ietf-spring-segment-routing] specifies that any instruction can be bound to a segment. Thus, an SR Policy can be built using any types of Segment Identifiers (SIDs) including those associated with topological or service instructions.

This section defines the key aspects and constituents of an SR Policy.

2.1. Identification of an SR Policy

An SR Policy is identified through the tuple <headend, color, endpoint>. In the context of a specific headend, one may identify an SR policy by the <color, endpoint> tuple.

The headend is the node where the policy is instantiated/implemented. The headend is specified as an IPv4 or IPv6 address.

The endpoint indicates the destination of the policy. The endpoint is specified as an IPv4 or IPv6 address. In a specific case (refer to Section 8.8.1), the endpoint can be the null address (0.0.0.0 for IPv4, ::0 for IPv6).

The color is a 32-bit numerical value that associates the SR Policy with an intent (e.g. low-latency).

The endpoint and the color are used to automate the steering of service or transport routes on SR Policies (refer to Section 8).

An implementation MAY allow assignment of a symbolic name comprising of printable ASCII characters to an SR Policy to serve as an user-friendly attribute for debug and troubleshooting purposes. Such symbolic names MUST NOT be considered as identifiers for an SR Policy.

2.2. Candidate Path and Segment List

An SR Policy is associated with one or more candidate paths. A candidate path is the unit for signaling of an SR Policy to a headend via protocols like Path Computation Element (PCE) Communication Protocol (PCEP) [RFC8281] or BGP SR Policy [I-D.ietf-idr-segment-routing-te-policy].

A candidate path is itself associated with a Segment-List (SID-List) or a set of SID-Lists.

A SID-List represents a specific source-routed way to send traffic from the headend to the endpoint of the corresponding SR policy.

A candidate path is either dynamic or explicit.

An explicit candidate path is associated with a SID-List or a set of SID-Lists.

A dynamic candidate path expresses an optimization objective and a set of constraints. The headend (potentially with the help of a PCE) computes the solution SID-List (or set of SID-Lists) that solves the optimization problem.

When a candidate path is associated with a set of SID-Lists, each SID-List is associated with a weight for weighted load balancing (refer Section 2.11 for details). The default weight is 1.

A variation of SR Policy can be used for packet replication. A candidate path could comprise multiple SID-Lists; one for each replication path. In such a scenario, packets are actually replicated through each SID List of the SR Policy to realize a point-to-multipoint service delivery. The weight of each SID-List does not come into the picture in this case since there is no load-balancing. The details of this and other such mechanisms for use of SR Policy for point-to-multipoint delivery are outside the scope of this document.

2.3. Protocol-Origin of a Candidate Path

A headend may be informed about a candidate path for an SR Policy <color, endpoint> by various means including: via configuration, PCEP [RFC8281] or BGP [I-D.ietf-idr-segment-routing-te-policy].

Protocol-Origin of a candidate path is an 8-bit value which identifies the component or protocol that originates or signals the candidate path.

The table below specifies the RECOMMENDED default values:

Value	Protocol-Origin
10	PCEP
20	BGP SR Policy
30	Local (via CLI, Yang model through NETCONF, gRPC, etc.)

Table 1: Protocol-origin Identifier

Implementations MAY allow modifications of these default values assigned to protocols on the headend along similar lines as a routing administrative distance. Its application in the candidate path selection is described in Section 2.9.

2.4. Originator of a Candidate Path

Originator identifies the node which provisioned or signalled the candidate path on the headend. The originator is expressed in the form of a 160 bit numerical value formed by the concatenation of the fields of the tuple <ASN, node-address> as below:

- ASN : represented as a 4 byte number.
- Node Address : represented as a 128 bit value. IPv4 addresses are encoded in the lowest 32 bits.

Its application in the candidate path selection is described in Section 2.9.

When Protocol-Origin is Local, the ASN and node address MAY be set to either the headend or the provisioning controller/node ASN and address. Default value is 0 for both AS and node address.

When Protocol-Origin is PCEP, it is the IPv4 or IPv6 address of the PCE and the AS number SHOULD be set to 0 by default when not available or known.

Protocol-Origin is BGP SR Policy, it is provided by the BGP component on the headend and is:

- the BGP Router ID and ASN of the node/controller signalling the candidate path when it has a BGP session to the headend, OR

- the BGP Router ID of the eBGP peer signalling the candidate path along with ASN of origin when the signalling is done via one or more intermediate eBGP routers, OR

- the BGP Originator ID [RFC4456] and the ASN of the node/controller when the signalling is done via one or more route-reflectors over iBGP session.

2.5. Discriminator of a Candidate Path

The Discriminator is a 32 bit value associated with a candidate path that uniquely identifies it within the context of an SR Policy from a specific Protocol-Origin as specified below:

When Protocol-Origin is Local, this is an implementation's configuration model specific unique identifier for a candidate path. Default value is 0.

When PCEP is the Protocol-Origin, the method to uniquely identify signalled path will be specified in a future PCEP document. Default value is 0.

When BGP SR Policy is the Protocol-Origin, it is the distinguisher specified in Section 2.1 of [I-D.ietf-idr-segment-routing-te-policy].

Its application in the candidate path selection is described in Section 2.9.

2.6. Identification of a Candidate Path

A candidate path is identified in the context of a single SR Policy.

A candidate path is not shared across SR Policies.

A candidate path is not identified by its SID-List(s).

> If CP1 is a candidate path of SR Policy Pol1 and CP2 is a candidate path of SR Policy Pol2, then these two candidate paths are independent, even if they happen to have the same SID-List. The SID-List does not identify a candidate path. The SID-List is an attribute of a candidate path.

The identity of a candidate path MUST be uniquely established in the context of an SR Policy <headend, color, endpoint> in order to handle add, delete or modify operations on them in an unambiguous manner regardless of their source(s).

The tuple <Protocol-Origin, originator, discriminator> uniquely identify a candidate path.

Candidate paths MAY also be assigned or signaled with a symbolic name comprising printable ASCII characters to serve as a user-friendly attribute for debug and troubleshooting purposes. Such symbolic names MUST NOT be considered as identifiers for a candidate path.

2.7. Preference of a Candidate Path

The preference of the candidate path is used to select the best candidate path for an SR Policy. The default preference is 100.

It is RECOMMENDED that each candidate path of a given SR policy has a different preference.

2.8. Validity of a Candidate Path

A candidate path is valid if it is usable. A common path validity criterion is the reachability of its constituent SIDs. The validation rules are specified in Section 5.

2.9. Active Candidate Path

A candidate path is selected when it is valid and it is determined to be the best path of the SR Policy. The selected path is referred to as the "active path" of the SR policy in this document.

Whenever a new path is learned or an active path is deleted, the validity of an existing path changes or an existing path is changed, the selection process MUST be re-executed.

The candidate path selection process operates on the candidate path Preference. A candidate path is selected when it is valid and it has the highest preference value among all the candidate paths of the SR Policy.

In the case of multiple valid candidate paths of the same preference, the tie-breaking rules are evaluated on the identification tuple in the following order until only one valid best path is selected:

1. Higher value of Protocol-Origin is selected.

2. Lower value of originator is selected.

3. Finally, the higher value of discriminator is selected.

An implementation MAY choose to override any of the tie-breaking rules above and maintain the already selected candidate path as active path.

The rules are framed with multiple protocols and sources in mind and hence may not follow the logic of a single protocol (e.g. BGP best path selection). The motivation behind these rules are as follows:

- The Protocol-Origin allows an operator to setup a default selection mechanism across protocol sources, e.g., to prefer locally provisioned over paths signalled via BGP SR Policy or PCEP.

- The preference, being the first tiebreaker, allows an operator to influence selection across paths thus allowing provisioning of multiple path options, e.g., CP1 is preferred and if it becomes invalid then fall-back to CP2 and so on. Since preference works across protocol sources it also enables (where necessary) selective override of the default protocol-origin preference, e.g., to prefer a path signalled via BGP SR Policy over what is locally provisioned.

- The originator allows an operator to have multiple redundant controllers and still maintain a deterministic behaviour over which of them are preferred even if they are providing the same candidate paths for the same SR policies to the headend.

- The discriminator performs the final tiebreaking step to ensure a deterministic outcome of selection regardless of the order in which candidate paths are signalled across multiple transport channels or sessions.

[I-D.filsfils-spring-sr-policy-considerations] provides a set of examples to illustrate the active candidate path selection rules.

2.10. Validity of an SR Policy

An SR Policy is valid when it has at least one valid candidate path.

2.11. Instantiation of an SR Policy in the Forwarding Plane

A valid SR Policy is instantiated in the forwarding plane.

Only the active candidate path is used for forwarding traffic that is being steered onto that policy.

If a set of SID-Lists is associated with the active path of the policy, then the steering is per flow and W-ECMP based according to the relative weight of each SID-List.

The fraction of the flows associated with a given SID-List is w/Sw where w is the weight of the SID-List and Sw is the sum of the weights of the SID-Lists of the selected path of the SR Policy.

The accuracy of the weighted load-balancing depends on the platform implementation.

2.12. Priority of an SR Policy

Upon topological change, many policies could be recomputed. An implementation MAY provide a per-policy priority configuration. The operator MAY set this field to indicate order in which the policies should be re-computed. Such a priority is represented by an integer in the range (0, 255) where the lowest value is the highest priority. The default value of priority is 128.

An SR Policy may comprise multiple Candidate Paths received from the same or different sources. A candidate path MAY be signaled with a priority value. When an SR Policy has multiple candidate paths with distinct signaled non-default priority values, the SR Policy as a whole takes the lowest value (i.e. the highest priority) amongst these signaled priority values.

2.13. Summary

In summary, the information model is the following:

```
         SR policy POL1 <headend, color, endpoint>
           Candidate-path CP1 <protocol-origin = 20, originator =
100:1.1.1.1, discriminator = 1>
             Preference 200
             Weight W1, SID-List1 <SID11...SID1i>
             Weight W2, SID-List2 <SID21...SID2j>
           Candidate-path CP2 <protocol-origin = 20, originator =
100:2.2.2.2, discriminator = 2>
             Preference 100
             Weight W3, SID-List3 <SID31...SID3i>
             Weight W4, SID-List4 <SID41...SID4j>
```

The SR Policy POL1 is identified by the tuple <headend, color, endpoint>. It has two candidate paths CP1 and CP2. Each is identified by a tuple <protocol-origin, originator, discriminator>. CP1 is the active candidate path (it is valid and it has the highest preference). The two SID-Lists of CP1 are installed as the

forwarding instantiation of SR policy Pol1. Traffic steered on Pol1 is flow-based hashed on SID-List <SID11...SID1i> with a ratio W1/(W1+W2).

3. Segment Routing Database

An SR headend maintains the Segment Routing Database (SR-DB).

An SR headend leverages the SR-DB to validate explicit candidate paths and compute dynamic candidate paths.

The information in the SR-DB MAY include:

- IGP information (topology, IGP metrics based on ISIS [RFC1195] and OSPF [RFC2328] [RFC5340])
- Segment Routing information (such as SRGB, SRLB, Prefix-SIDs, Adj-SIDs, BGP Peering SID, SRv6 SIDs) [I-D.ietf-spring-segment-routing] [I-D.ietf-idr-bgpls-segment-routing-epe] [I-D.filsfils-spring-srv6-network-programming]
- TE Link Attributes (such as TE metric, SRLG, attribute-flag, extended admin group) [RFC5305] [RFC3630].
- Extended TE Link attributes (such as latency, loss) [RFC7810] [RFC7471]
- Inter-AS Topology information [I-D.ietf-idr-bgpls-segment-routing-epe].

The attached domain topology MAY be learned via IGP, BGP-LS or NETCONF.

A non-attached (remote) domain topology MAY be learned via BGP-LS or NETCONF.

In some use-cases, the SR-DB may only contain the attached domain topology while in others, the SR-DB may contain the topology of multiple domains and in this case it is multi-domain capable.

The SR-DB MAY also contain the SR Policies instantiated in the network. This can be collected via BGP-LS [I-D.ietf-idr-te-lsp-distribution] or PCEP [RFC8231] and [I-D.sivabalan-pce-binding-label-sid]. This information allows to build an end-to-end policy on the basis of intermediate SR policies (see Section 6 for further details).

The SR-DB MAY also contain the Maximum SID Depth (MSD) capability of nodes in the topology. This can be collected via ISIS [I-D.ietf-isis-segment-routing-msd], OSPF [I-D.ietf-ospf-segment-routing-msd], BGP-LS

[I-D.ietf-idr-bgp-ls-segment-routing-msd] or PCEP [I-D.ietf-pce-segment-routing].

The use of the SR-DB for computation and validation of SR Policies is outside the scope of this document. Some implementation aspects related to this are covered in [I-D.filsfils-spring-sr-policy-considerations].

4. Segment Types

A SID-List is an ordered set of segments represented as <S1, S2, ... Sn> where S1 is the first segment.

Based on the desired dataplane, either the MPLS label stack or the SRv6 SRH is built from the SID-List. However, the SID-List itself can specified using different segment-descriptor types and the following are currently defined:

Type 1: SR-MPLS Label:
A MPLS label corresponding to any of the segment types defined for SR-MPLS (as defined in [I-D.ietf-spring-segment-routing] or other SR-MPLS specifications) can be used. Additionally, reserved labels like explicit-null or in general any MPLS label may also be used. e.g. this type can be used to specify a label representation which maps to an optical transport path on a packet transport node. This type does not require the headend to perform SID resolution.

Type 2: SRv6 SID:
An IPv6 address corresponding to any of the segment types defined for SRv6 (as defined in [I-D.filsfils-spring-srv6-network-programming] or other SRv6 specifications) can be used. This type does not require the headend to perform SID resolution.

Type 3: IPv4 Prefix with optional SR Algorithm:
The headend is required to resolve the specified IPv4 Prefix Address to the SR-MPLS label corresponding to a Prefix SID segment (as defined in [I-D.ietf-spring-segment-routing]). The SR algorithm (refer to Section 3.1.1 of [I-D.ietf-spring-segment-routing]) to be used MAY also be provided.

Type 4: IPv6 Global Prefix with optional SR Algorithm for SR-MPLS:
In this case the headend is required to resolve the specified IPv6 Global Prefix Address to the SR-MPLS label corresponding to its Prefix SID segment (as defined in [I-D.ietf-spring-segment-routing]). The SR Algorithm (refer to

Section 3.1.1 of [I-D.ietf-spring-segment-routing]) to be used MAY also be provided.

Type 5: IPv4 Prefix with Local Interface ID:
This type allows identification of Adjacency SID (as defined in [I-D.ietf-spring-segment-routing]) or BGP EPE Peering SID (as defined in [I-D.ietf-idr-bgpls-segment-routing-epe]) label for point-to-point links including IP unnumbered links. The headend is required to resolve the specified IPv4 Prefix Address to the Node originating it and then use the Local Interface ID to identify the point-to-point link whose adjacency is being referred to. The Local Interface ID link descriptor follows semantics as specified in [RFC7752]. This type can also be used to indicate indirection into a layer 2 interface (i.e. without IP address) like a representation of an optical transport path or a layer 2 Ethernet port or circuit at the specified node.

Type 6: IPv4 Addresses for link endpoints as Local, Remote pair:
This type allows identification of Adjacency SID (as defined in [I-D.ietf-spring-segment-routing]) or BGP EPE Peering SID (as defined in [I-D.ietf-idr-bgpls-segment-routing-epe]) label for links. The headend is required to resolve the specified IPv4 Local Address to the Node originating it and then use the IPv4 Remote Address to identify the link adjacency being referred to. The Local and Remote Address pair link descriptors follows semantics as specified in [RFC7752].

Type 7: IPv6 Prefix and Interface ID for link endpoints as Local, Remote pair for SR-MPLS:
This type allows identification of Adjacency SID (as defined in [I-D.ietf-spring-segment-routing]) or BGP EPE Peering SID (as defined in [I-D.ietf-idr-bgpls-segment-routing-epe]) label for links including those with only Link Local IPv6 addresses. The headend is required to resolve the specified IPv6 Prefix Address to the Node originating it and then use the Local Interface ID to identify the point-to-point link whose adjacency is being referred to. For other than point-to-point links, additionally the specific adjacency over the link needs to be resolved using the Remote Prefix and Interface ID. The Local and Remote pair of Prefix and Interface ID link descriptor follows semantics as specified in [RFC7752]. This type can also be used to indicate indirection into a layer 2 interface (i.e. without IP address) like a representation of an optical transport path or a layer 2 Ethernet port or circuit at the specified node.

Type 8: IPv6 Addresses for link endpoints as Local, Remote pair for SR-MPLS:

This type allows identification of Adjacency SID (as defined in [I-D.ietf-spring-segment-routing]) or BGP EPE Peering SID (as defined in [I-D.ietf-idr-bgpls-segment-routing-epe]) label for links with Global IPv6 addresses. The headend is required to resolve the specified Local IPv6 Address to the Node originating it and then use the Remote IPv6 Address to identify the link adjacency being referred to. The Local and Remote Address pair link descriptors follows semantics as specified in [RFC7752].

Type 9: IPv6 Global Prefix with optional SR Algorithm for SRv6:

The headend is required to resolve the specified IPv6 Global Prefix Address to the SRv6 END function SID (as defined in [I-D.filsfils-spring-srv6-network-programming]) corresponding to the node which is originating the prefix. The SR Algorithm (refer to Section 3.1.1 of [I-D.ietf-spring-segment-routing]) to be used MAY also be provided.

Type 10: IPv6 Prefix and Interface ID for link endpoints as Local, Remote pair for SRv6:

This type allows identification of SRv6 END.X SID (as defined in [I-D.filsfils-spring-srv6-network-programming]) for links with only Link Local IPv6 addresses. The headend is required to resolve the specified IPv6 Prefix Address to the Node originating it and then use the Local Interface ID to identify the point-to-point link whose adjacency is being referred to. For other than point-to-point links, additionally the specific adjacency needs to be resolved using the Remote Prefix and Interface ID. The Local and Remote pair of Prefix and Interface ID link descriptor follows semantics as specified in [RFC7752].

Type 11: IPv6 Addresses for link endpoints as Local, Remote pair for SRv6:

This type allows identification of SRv6 END.X SID (as defined in [I-D.filsfils-spring-srv6-network-programming]) for links with Global IPv6 addresses. The headend is required to resolve the specified Local IPv6 Address to the Node originating it and then use the Remote IPv6 Address to identify the link adjacency being referred to. The Local and Remote Address pair link descriptors follows semantics as specified in [RFC7752].

When the algorithm is not specified for the SID types above which optionally allow for it, the headend SHOULD use the Strict Shortest Path algorithm if available; otherwise it SHOULD use the default

Shortest Path algorithm. The specification of algorithm enables the use of IGP Flex Algorithm [I-D.ietf-lsr-flex-algo] specific SIDs in SR Policy.

For SID types 3-through-11, a SID value may also be optionally provided to the headend for verification purposes. Section 5.1. describes the resolution and verification of the SIDs and Segment Lists on the headend.

When building the MPLS label stack or the IPv6 Segment list from the Segment List, the node instantiating the policy MUST interpret the set of Segments as follows:

- The first Segment represents the topmost label or the first IPv6 segment. It identifies the first segment the traffic will be directed toward along the SR explicit path.
- The last Segment represents the bottommost label or the last IPv6 segment the traffic will be directed toward along the SR explicit path.

4.1. Explicit Null

A Type 1 SID may be any MPLS label, including reserved labels.

For example, assuming that the desired traffic-engineered path from a headend 1 to an endpoint 4 can be expressed by the SID-List <16002, 16003, 16004> where 16002, 16003 and 16004 respectively refer to the IPv4 Prefix SIDs bound to node 2, 3 and 4, then IPv6 traffic can be traffic-engineered from nodes 1 to 4 via the previously described path using an SR Policy with SID-List <16002, 16003, 16004, 2> where mpls label value of 2 represents the "IPv6 Explicit NULL Label".

The penultimate node before node 4 will pop 16004 and will forward the frame on its directly connected interface to node 4.

The endpoint receives the traffic with top label "2" which indicates that the payload is an IPv6 packet.

When steering unlabeled IPv6 BGP destination traffic using an SR policy composed of SID-List(s) based on IPv4 SIDs, the Explicit Null Label Policy is processed as specified in [I-D.ietf-idr-segment-routing-te-policy]) Section 2.4.4. When an "IPv6 Explicit NULL label" is not present as the bottom label, the headend SHOULD automatically impose one. Refer to Section 8) later in this document for more details.

5. Validity of a Candidate Path

5.1. Explicit Candidate Path

An explicit candidate path is associated with a SID-List or a set of SID-Lists.

An explicit candidate path is provisioned by the operator directly or via a controller.

The computation/logic that leads to the choice of the SID list is external to the SR Policy headend. The SR Policy headend does not compute the SID list. The SR Policy headend only confirms its validity.

A SID-List of an explicit candidate path MUST be declared invalid when:

- It is empty.
- Its weight is 0.
- The headend is unable to resolve the first SID into one or more outgoing interface(s) and next-hop(s).
- The headend is unable to resolve any non-first SID of type 3-through-11 into an MPLS label or an SRv6 SID.
- The headend verification fails for any SID for which verification has been explicitly requested.

"Unable to resolve" means that the headend has no path to the SID in its SR database.

SID verification is performed when the headend is explicitly requested to verify SID(s) by the controller via the signaling protocol used. Implementations MAY provide a local configuration option to enable verification on a global or per policy or per candidate path basis.

"Verification fails" for a SID means any of the following:

- The headend is unable to find the SID in its SR DB
- The headend detects mis-match between the SID value and its context provided for SIDs of type 3-through-11 in its SR DB.
- The headend is unable to resolve any non-first SID of type 3-through-11 into an MPLS label or an SRv6 SID.

In multi-domain deployments, it is expected that the headend be unable to verify the reachability of the SIDs in remote domains. Types A or B MUST be used for the SIDs for which the reachability

cannot be verified. Note that the first SID MUST always be reachable regardless of its type.

In addition, a SID-List MAY be declared invalid when:

- Its last segment is not a Prefix SID (including BGP Peer Node-SID) advertised by the node specified as the endpoint of the corresponding SR policy.
- Its last segment is not an Adjacency SID (including BGP Peer Adjacency SID) of any of the links present on neighbor nodes and that terminate on the node specified as the endpoint of the corresponding SR policy.

An explicit candidate path is invalid as soon as it has no valid SID-List.

5.2. Dynamic Candidate Path

A dynamic candidate path is specified as an optimization objective and constraints.

The headend of the policy leverages its SR database to compute a SID-List ("solution SID-List") that solves this optimization problem.

The headend re-computes the solution SID-List any time the inputs to the problem change (e.g., topology changes).

When local computation is not possible (e.g., a policy's tailend is outside the topology known to the headend) or not desired, the headend MAY send path computation request to a PCE supporting PCEP extension specified in [I-D.ietf-pce-segment-routing].

If no solution is found to the optimization objective and constraints, then the dynamic candidate path MUST be declared invalid.

[I-D.filsfils-spring-sr-policy-considerations] discusses some of the optimization objectives and constraints that may be considered by a dynamic candidate path. It illustrates some of the desirable properties of the computation of the solution SID list.

6. Binding SID

The Binding SID (BSID) is fundamental to Segment Routing [I-D.ietf-spring-segment-routing]. It provides scaling, network opacity and service independence. [I-D.filsfils-spring-sr-policy-considerations] illustrates some of these benefits.

6.1. BSID of a candidate path

Each candidate path MAY be defined with a BSID.

Candidate Paths of the same SR policy SHOULD have the same BSID.

Candidate Paths of different SR policies MUST NOT have the same BSID.

6.2. BSID of an SR Policy

The BSID of an SR policy is the BSID of its active candidate path.

When the active candidate path has a specified BSID, the SR Policy uses that BSID if this value (label in MPLS, IPv6 address in SRv6) is available (i.e., not associated with any other usage: e.g. to another MPLS client, to another SID, to another SR Policy).

Optionally, instead of only checking that the BSID of the active path is available, a headend MAY check that it is available within a given SID range i.e., Segment Routing Local Block (SRLB) as specified in [I-D.ietf-spring-segment-routing].

When the specified BSID is not available (optionally is not in the SRLB), an alert message is generated.

In the cases (as described above) where SR Policy does not have a BSID available, then the SR Policy MAY dynamically bind a BSID to itself. Dynamically bound BSID SHOULD use an available SID outside the SRLB.

Assuming that at time t the BSID of the SR Policy is B1, if at time t+dt a different candidate path becomes active and this new active path does not have a specified BSID or its BSID is specified but is not available, then the SR Policy keeps the previous BSID B1.

The association of an SR Policy with a BSID thus MAY change over the life of the SR policy (e.g., upon active path change). Hence, the BSID cannot be used as an identification of an SR Policy.

6.2.1. Frequent use-cases : unspecified BSID

All the candidate paths of the same SR Policy can have an unspecified BSID.

In such a case, a BSID MAY be dynamically bound to the SR Policy as soon as the first valid candidate path is received. That BSID is kept along all the life of the SR Policy and across changes of active candidate path.

6.2.2. Frequent use-case: all specified to the same BSID

All the paths of the SR Policy can have the same specified BSID.

6.2.3. Specified-BSID-only

An implementation MAY support the configuration of the Specified-BSID-only restrictive behavior on the headend for all SR Policies or individual SR Policies. Further, this restrictive behavior MAY also be signaled on a per SR Policy basis to the headend.

When this restrictive behavior is enabled, if the candidate path has an unspecified BSID or if the specified BSID is not available when the candidate path becomes active then no BSID is bound to it and it is considered invalid. An alert is triggered. Other candidate paths can then be evaluated for becoming the active candidate path.

6.3. Forwarding Plane

A valid SR Policy installs a BSID-keyed entry in the forwarding plane with the action of steering the packets matching this entry to the selected path of the SR Policy.

If the Specified-BSID-only restrictive behavior is enabled and the BSID of the active path is not available (optionally not in the SRLB), then the SR Policy does not install any entry indexed by a BSID in the forwarding plane.

6.4. Non-SR usage of Binding SID

An implementation MAY choose to associate a Binding SID with any type of interface (e.g. a layer 3 termination of an Optical Circuit) or a tunnel (e.g. IP tunnel, GRE tunnel, IP/UDP tunnel, MPLS RSVP-TE tunnel, etc). This enables the use of other non-SR enabled interfaces and tunnels as segments in an SR Policy SID-List without the need of forming routing protocol adjacencies over them.

The details of this kind of usage are beyond the scope of this document. A specific packet optical integration use case is described in [I-D.anand-spring-poi-sr]

7. SR Policy State

The SR Policy State is maintained on the headend to represent the state of the policy and its candidate paths. This is to provide an accurate representation of whether the SR Policy is being instantiated in the forwarding plane and which of its candidate paths and segment-list(s) are active. The SR Policy state MUST also

reflect the reason when a policy and/or its candidate path is not active due to validation errors or not being preferred.

The SR Policy state can be reported by the headend node via BGP-LS [I-D.ietf-idr-te-lsp-distribution] or PCEP [RFC8231] and [I-D.sivabalan-pce-binding-label-sid].

SR Policy state on the headend also includes traffic accounting information for the flows being steered via the policies. The details of the SR Policy accounting are beyond the scope of this document and [I-D.ali-spring-sr-traffic-accounting] covers these aspects in the broader context of traffic accounting in a SR network.

Implementations MAY support an administrative state to control locally provisioned policies via mechanisms like CLI or NETCONF.

8. Steering into an SR Policy

A headend can steer a packet flow into a valid SR Policy in various ways:

- Incoming packets have an active SID matching a local BSID at the headend.
- Per-destination Steering: incoming packets match a BGP/Service route which recurses on an SR policy.
- Per-flow Steering: incoming packets match or recurse on a forwarding array of where some of the entries are SR Policies.
- Policy-based Steering: incoming packets match a routing policy which directs them on an SR policy.

For simplicity of illustration, this document uses the SR-MPLS example.

8.1. Validity of an SR Policy

An SR Policy is invalid when all its candidate paths are invalid.

By default, upon transitioning to the invalid state,

- an SR Policy and its BSID are removed from the forwarding plane.
- any steering of a service (PW), destination (BGP-VPN), flow or packet on the related SR policy is disabled and the related service, destination, flow or packet is routed per the classic forwarding table (e.g. longest-match to the destination or the recursing next-hop).

8.2. Drop upon invalid SR Policy

An SR Policy MAY be enabled for the Drop-Upon-Invalid behavior:

- an invalid SR Policy and its BSID is kept in the forwarding plane with an action to drop.
- any steering of a service (PW), destination (BGP-VPN), flow or packet on the related SR policy is maintained with the action to drop all of this traffic.

The drop-upon-invalid behavior has been deployed in use-cases where the operator wants some PW to only be transported on a path with specific constraints. When these constraints are no longer met, the operator wants the PW traffic to be dropped. Specifically, the operator does not want the PW to be routed according to the IGP shortest-path to the PW endpoint.

8.3. Incoming Active SID is a BSID

Let us assume that headend H has a valid SR Policy P of SID-List <S1, S2, S3> and BSID B.

When H receives a packet K with label stack <B, L2, L3>, H pops B and pushes <S1, S2, S3> and forwards the resulting packet according to SID S1.

> "Forwarding the resulting packet according to S1" means: If S1 is an Adj SID or a PHP-enabled prefix SID advertised by a neighbor, H sends the resulting packet with label stack <S2, S3, L2, L3> on the outgoing interface associated with S1; Else H sends the resulting packet with label stack <S1, S2, S3, L2, L3> along the path of S1.

H has steered the packet in the SR policy P.

H did not have to classify the packet. The classification was done by a node upstream of H (e.g., the source of the packet or an intermediate ingress edge node of the SR domain) and the result of this classification was efficiently encoded in the packet header as a BSID.

This is another key benefit of the segment routing in general and the binding SID in particular: the ability to encode a classification and the resulting steering in the packet header to better scale and simplify intermediate aggregation nodes.

If the SR Policy P is invalid, the BSID B is not in the forwarding plane and hence the packet K is dropped by H.

8.4. Per-Destination Steering

Let us assume that headend H:

- learns a BGP route R/r via next-hop N, extended-color community C and VPN label V.
- has a valid SR Policy P to (endpoint = N, color = C) of SID-List <S1, S2, S3> and BSID B.
- has a BGP policy which matches on the extended-color community C and allows its usage as SLA steering information.

If all these conditions are met, H installs R/r in RIB/FIB with next-hop = SR Policy P of BSID B instead of via N.

Indeed, H's local BGP policy and the received BGP route indicate that the headend should associate R/r with an SR Policy path to N with the SLA associated with color C. The headend therefore installs the BGP route on that policy.

This can be implemented by using the BSID as a generalized next-hop and installing the BGP route on that generalized next-hop.

When H receives a packet K with a destination matching R/r, H pushes the label stack <S1, S2, S3, V> and sends the resulting packet along the path to S1.

Note that any SID associated with the BGP route is inserted after the SID-List of the SR Policy (i.e., <S1, S2, S3, V>).

The same behavior is applicable to any type of service route: any AFI/SAFI of BGP [RFC4760] any AFI/SAFI of LISP [RFC6830].

8.4.1. Multiple Colors

When a BGP route has multiple extended-color communities each with a valid SR Policy NLRI, the BGP process installs the route on the SR policy whose color is of highest numerical value.

Let us assume that headend H:

- learns a BGP route R/r via next-hop N, extended-color communities C1 and C2 and VPN label V.
- has a valid SR Policy P1 to (endpoint = N, color = C1) of SID list <S1, S2, S3> and BSID B1.
- has a valid SR Policy P2 to (endpoint = N, color = C2) of SID list <S4, S5, S6> and BSID B2.
- has a BGP policy which matches on the extended-color communities C1 and C2 and allows their usage as SLA steering information

If all these conditions are met, H installs R/r in RIB/FIB with next-hop = SR Policy P2 of BSID=B2 (instead of N) because C2 > C1.

8.5. Recursion on an on-demand dynamic BSID

In the previous section, it was assumed that H had a pre-established "explicit" SR Policy (endpoint N, color C).

In this section, independently to the a-priori existence of any explicit candidate path of the SR policy (N, C), it is to be noted that the BGP process at headend node H triggers the instantiation of a dynamic candidate path for the SR policy (N, C) as soon as:

- the BGP process learns of a route R/r via N and with color C.
- a local policy at node H authorizes the on-demand SR Policy path instantiation and maps the color to a dynamic SR Policy path optimization template.

8.5.1. Multiple Colors

When a BGP route R/r via N has multiple extended-color communities Ci (with i=1 ... n), an individual on-demand SR Policy dynamic path request (endpoint N, color Ci) is triggered for each color Ci.

8.6. Per-Flow Steering

Let us assume that headend H:

- has a valid SR Policy P1 to (endpoint = N, color = C1) of SID-List <S1, S2, S3> and BSID B1.
- has a valid SR Policy P2 to (endpoint = N, color = C2) of SID-List <S4, S5, S6> and BSID B2.
- is configured to instantiate an array of paths to N where the entry 0 is the IGP path to N, color C1 is the first entry and Color C2 is the second entry. The index into the array is called a Forwarding Class (FC). The index can have values 0 to 7.
- is configured to match flows in its ingress interfaces (upon any field such as Ethernet destination/source/vlan/tos or IP destination/source/DSCP or transport ports etc.) and color them with an internal per-packet forwarding-class variable (0, 1 or 2 in this example).

If all these conditions are met, H installs in RIB/FIB:

- N via a recursion on an array A (instead of the immediate outgoing link associated with the IGP shortest-path to N).
- Entry A(0) set to the immediate outgoing link of the IGP shortest-path to N.

- Entry A(1) set to SR Policy P1 of BSID=B1.
- Entry A(2) set to SR Policy P2 of BSID=B2.

H receives three packets K, K1 and K2 on its incoming interface. These three packets either longest-match on N or more likely on a BGP/service route which recurses on N. H colors these 3 packets respectively with forwarding-class 0, 1 and 2. As a result:

- H forwards K along the shortest-path to N (which in SR-MPLS results in the pushing of the prefix-SID of N).
- H pushes <S1, S2, S3> on packet K1 and forwards the resulting frame along the shortest-path to S1.
- H pushes <S4, S5, S6> on packet K2 and forwards the resulting frame along the shortest-path to S4.

If the local configuration does not specify any explicit forwarding information for an entry of the array, then this entry is filled with the same information as entry 0 (i.e. the IGP shortest-path).

If the SR Policy mapped to an entry of the array becomes invalid, then this entry is filled with the same information as entry 0. When all the array entries have the same information as entry0, the forwarding entry for N is updated to bypass the array and point directly to its outgoing interface and next-hop.

The array index values (e.g. 0, 1 and 2) and the notion of forwarding-class are implementation specific and only meant to describe the desired behavior. The same can be realized by other mechanisms.

This realizes per-flow steering: different flows bound to the same BGP endpoint are steered on different IGP or SR Policy paths.

8.7. Policy-based Routing

Finally, headend H may be configured with a local routing policy which overrides any BGP/IGP path and steer a specified packet on an SR Policy. This includes the use of mechanisms like IGP Shortcut for automatic routing of IGP prefixes over SR Policies intended for such purpose.

8.8. Optional Steering Modes for BGP Destinations

8.8.1. Color-Only BGP Destination Steering

In the previous section, it is seen that the steering on an SR Policy is governed by the matching of the BGP route's next-hop N and the authorized color C with an SR Policy defined by the tuple (N, C).

This is the most likely form of BGP destination steering and the one recommended for most use-cases.

This section defines an alternative steering mechanism based only on the color.

This color-only steering variation is governed by two new flags "C" and "O" defined in the color extended community [ref draft-ietf-idr-segment-routing-te-policy section 3].

The Color-Only flags "CO" are set to 00 by default.

When 00, the BGP destination is steered as follows:

```
    IF there is a valid SR Policy (N, C) where N is the IPv4 or IPv6

            endpoint address and C is a color;
        Steer into SR Policy (N, C);
    ELSE;
        Steer on the IGP path to the next-hop N.
```

This is the classic case described in this document previously and what is recommended in most scenarios.

When 01, the BGP destination is steered as follows:

```
    IF there is a valid SR Policy (N, C) where N is the IPv4 or IPv6

            endpoint address and C is a color;
        Steer into SR Policy (N, C);
    ELSE IF there is a valid SR Policy (null endpoint, C) of the
            same address-family of N;
        Steer into SR Policy (null endpoint, C);
    ELSE IF there is any valid SR Policy
            (any address-family null endpoint, C);
        Steer into SR Policy (any null endpoint, C);
    ELSE;
        Steer on the IGP path to the next-hop N.
```

When 10, the BGP destination is steered as follows:

```
    IF there is a valid SR Policy (N, C) where N is an IPv4 or IPv6
            endpoint address and C is a color;
        Steer into SR Policy (N, C);
    ELSE IF there is a valid SR Policy (null endpoint, C)
            of the same address-family of N;
        Steer into SR Policy (null endpoint, C);
    ELSE IF there is any valid SR Policy
```

```
            (any address-family null endpoint, C);
        Steer into SR Policy (any null endpoint, C);
    ELSE IF there is any valid SR Policy (any endpoint, C)
            of the same address-family of N;
        Steer into SR Policy (any endpoint, C);
    ELSE IF there is any valid SR Policy
            (any address-family endpoint, C);
        Steer into SR Policy (any address-family endpoint, C);
    ELSE;
        Steer on the IGP path to the next-hop N.
```

The null endpoint is 0.0.0.0 for IPv4 and ::0 for IPv6 (all bits set to the 0 value).

The value 11 is reserved for future use and SHOULD NOT be used. Upon reception, an implementations MUST treat it like 00.

8.8.2. Multiple Colors and CO flags

The steering preference is first based on highest color value and then CO-dependent for the color. Assuming a Prefix via (NH, C1(CO=01), C2(CO=01)); C1>C2 The steering preference order is:

- SR policy (NH, C1).
- SR policy (null, C1).
- SR policy (NH, C2).
- SR policy (null, C2).
- IGP to NH.

8.8.3. Drop upon Invalid

This document defined earlier that when all the following conditions are met, H installs R/r in RIB/FIB with next-hop = SR Policy P of BSID B instead of via N.

- H learns a BGP route R/r via next-hop N, extended-color community C and VPN label V.
- H has a valid SR Policy P to (endpoint = N, color = C) of SID-List <S1, S2, S3> and BSID B.
- H has a BGP policy which matches on the extended-color community C and allows its usage as SLA steering information.

This behavior is extended by noting that the BGP policy may require the BGP steering to always stay on the SR policy whatever its validity.

This is the "drop upon invalid" option described in section 10.2 applied to BGP-based steering.

9. Protection

9.1. Leveraging TI-LFA local protection of the constituent IGP segments

In any topology, Topology-Independent Loop Free Alternate (TI-LFA) [I-D.bashandy-rtgwg-segment-routing-ti-lfa] provides a 50msec local protection technique for IGP SIDs. The backup path is computed on a per IGP SID basis along the post-convergence path.

In a network that has deployed TI-LFA, an SR Policy built on the basis of TI-LFA protected IGP segments leverages the local protection of the constituent segments.

In a network that has deployed TI-LFA, an SR Policy instantiated only with non-protected Adj SIDs does not benefit from any local protection.

9.2. Using an SR Policy to locally protect a link

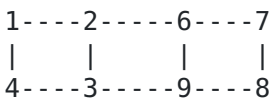

Figure 1: Local protection using SR Policy

An SR Policy can be instantiated at node 2 to protect the link 2to6. A typical explicit SID list would be <3, 9, 6>.

A typical use-case occurs for links outside an IGP domain: e.g. 1, 2, 3 and 4 are part of IGP/SR sub-domain 1 while 6, 7, 8 and 9 are part of IGP/SR sub-domain 2. In such a case, links 2to6 and 3to9 cannot benefit from TI-LFA automated local protection.

9.3. Using a Candidate Path for Path Protection

An SR Policy allows for multiple candidate paths, of which at any point in time there is a single active candidate path that is provisioned in the forwarding plane and used for traffic steering. However, another (lower preference) candidate path MAY be designated as the backup for a specific or all (active) candidate path(s). Such a backup candidate path is generally disjoint from the active candidate path.

The headend MAY compute a-priori and validate such backup candidate paths as well as provision them into forwarding plane as backup for the active path. A fast re-route mechanism MAY then be used to

trigger sub 50msec switchover from the active to the backup candidate path in the forwarding plane. Mechanisms like BFD MAY be used for fast detection of such failures.

10. Security Considerations

This document does not define any new protocol extensions and does not impose any additional security challenges.

11. IANA Considerations

This document has no actions for IANA.

12. Acknowledgement

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