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GOST R 34.12-2015: Block Cipher "Kuznyechik"
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Abstract

This document is intended to be a source of information about the Russian Federal standard block cipher with block length of $n=128$ bits, which is also referred as "Kuznyechik" [[GOST3412-2015](#)]. This algorithm is one of the Russian cryptographic standard algorithms (called GOST algorithms).

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1. Scope

The Russian Federal standard [[GOST3412-2015](#)] specifies basic block ciphers used as cryptographic techniques for information processing and information protection including the provision of confidentiality, authenticity, and integrity of information during information transmission, processing and storage in computer-aided systems.

The cryptographic algorithms specified in this Standard are designed both for hardware and software implementation. They comply with modern cryptographic requirements, and put no restrictions on the confidentiality level of the protected information.

The Standard applies to developing, operation, and modernization of the information systems of various purposes.

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [[RFC2119](#)].

2. General Information

The block cipher "Kuznyechik" [[GOST3412-2015](#)] was developed by the Center for Information Protection and Special Communications of the Federal Security Service of the Russian Federation with participation of the Open Joint-Stock company "Information Technologies and Communication Systems" (InfoTeCS JSC). GOST R 34.12-2015 was approved and introduced by Decree #749 of the Federal Agency on Technical Regulating and Metrology on 19.06.2015.

Terms and concepts in the standard comply with the following international standards:

- o ISO/IEC 10116 [[ISO-IEC10116](#)],
- o series of standards ISO/IEC 18033 [[ISO-IEC18033-1](#)], [[ISO-IEC18033-3](#)].

3. Definitions and Notations

The following terms and their corresponding definitions are used in the standard.

3.1. Definitions

Definitions

encryption algorithm: process which transforms plaintext into ciphertext (Clause 2.19 of [[ISO-IEC18033-1](#)]),

decryption algorithm: process which transforms ciphertext into plaintext (Clause 2.14 of [[ISO-IEC18033-1](#)]),

basic block cipher: block cipher which for a given key provides a single invertible mapping of the set of fixed-length plaintext blocks into ciphertext blocks of the same length,

block: string of bits of a defined length (Clause 2.6 of [[ISO-IEC18033-1](#)]),

block cipher: symmetric encipherment system with the property that the encryption algorithm operates on a block of plaintext, i.e. a string of bits of a defined length, to yield a block of ciphertext (Clause 2.7 of [[ISO-IEC18033-1](#)]),

Note: In GOST R 34.12-2015, it is established that the terms "block cipher" and "block encryption algorithm" are synonyms.

encryption: reversible transformation of data by a cryptographic algorithm to produce ciphertext, i.e., to hide the information content of the data (Clause 2.18 of [\[ISO-IEC18033-1\]](#)),

round key: sequence of symbols which is calculated from the key and controls a transformation for one round of a block cipher,

key: sequence of symbols that controls the operation of a cryptographic transformation (e.g., encipherment, decipherment) (Clause 2.21 of [\[ISO-IEC18033-1\]](#)),

Note: In GOST R 34.12-2015, the key must be a binary sequence.

plaintext: unencrypted information (Clause 3.11 of [\[ISO-IEC10116\]](#)),

key schedule: calculation of round keys from the key,

decryption: reversal of a corresponding encipherment (Clause 2.13 of [\[ISO-IEC18033-1\]](#)),

symmetric cryptographic technique: cryptographic technique that uses the same secret key for both the originator's and the recipient's transformation (Clause 2.32 of [\[ISO-IEC18033-1\]](#)),

cipher: alternative term for encipherment system (Clause 2.20 of [\[ISO-IEC18033-1\]](#)),

ciphertext: data which has been transformed to hide its information content (Clause 3.3 of [\[ISO-IEC10116\]](#)).

3.2. Notations

The following notations are used in the standard:

V^* the set of all binary vector-strings of a finite length (hereinafter referred to as the strings) including empty string,

V_s the set of all binary strings of length s , where s is a non-negative integer; substrings and string components are enumerated from right to left starting from zero,

$U[*]W$ direct (Cartesian) product of two set U s and W ,

$|A|$ the number of components (the length) of a string A belonging to V^* (if A is an empty string, then $|A| = 0$),

$A||B$ concatenation of strings A , B both belonging to V^* , i.e., a string from $V_{(|A|+|B|)}$, where the left substring from $V_{|A|}$ is equal to A and the right substring from $V_{|B|}$ is equal to B ,

$Z_{(2^n)}$ ring of residues modulo 2^n ,

Q finite field $GF(2)[x]/p(x)$, where $p(x)=x^8+x^7+x^6+x+1$ belongs to $GF(2)[x]$; elements of field Q are represented by integers in such way that element $z_0+z_1\theta+\dots+z_7\theta^7$ belonging to Q corresponds to integer $z_0+2z_1+\dots+2^7z_7$, where $z_i=0$ or $z_i=1$, $i=0,1,\dots,7$ and θ denotes a residue class modulo $p(x)$ containing x ,

(xor) exclusive-or of the two binary strings of the same length,

$Vec_s: Z_{(2^s)} \rightarrow V_s$ bijective mapping which maps an element from ring $Z_{(2^s)}$ into its binary representation, i.e., for an element z of the ring $Z_{(2^s)}$, represented by the residue $z_0 + (2z_1) + \dots + (2^{s-1}z_{s-1})$, where $z_i \in \{0, 1\}$, $i = 0, \dots, s-1$, the equality $Vec_s(z) = z_{s-1}||\dots||z_1||z_0$ holds,

$Int_s: V_s \rightarrow Z_{(2^s)}$ the mapping inverse to the mapping Vec_s , i.e., $Int_s = Vec_s^{-1}$,

$\nabla_8: V_8 \rightarrow Q$ bijective mapping which maps a binary string from V_8 into an element from field Q as follows: string $z_7||\dots||z_1||z_0$, where $z_i \in \{0, 1\}$, $i = 0, \dots, 7$, corresponds to the element $z_0+(z_1\theta)+\dots+(z_7\theta^7)$ belonging to Q ,

$\Delta: Q \rightarrow V_8$ the mapping inverse to the mapping ∇_8 , i.e., $\Delta = \nabla_8^{-1}$,

PS composition of mappings, where the mapping S applies first,

P^s composition of mappings P^{s-1} and P , where $P^1=P$,

4. Parameter Values

4.1. Nonlinear Bijection

The bijective nonlinear mapping is a substitution: $\Pi = (Vec_8)\Pi'(Int_8): V_8 \rightarrow V_8$, where $\Pi': Z_{(2^8)} \rightarrow Z_{(2^8)}$. The values of the substitution Π' are specified below as an array $\Pi' = (\Pi'(0), \Pi'(1), \dots, \Pi'(255))$:

Pi' =
 (252, 238, 221, 17, 207, 110, 49, 22, 251, 196, 250,
 218, 35, 197, 4, 77, 233, 119, 240, 219, 147, 46,
 153, 186, 23, 54, 241, 187, 20, 205, 95, 193, 249,
 24, 101, 90, 226, 92, 239, 33, 129, 28, 60, 66,
 139, 1, 142, 79, 5, 132, 2, 174, 227, 106, 143,
 160, 6, 11, 237, 152, 127, 212, 211, 31, 235, 52,
 44, 81, 234, 200, 72, 171, 242, 42, 104, 162, 253,
 58, 206, 204, 181, 112, 14, 86, 8, 12, 118, 18,
 191, 114, 19, 71, 156, 183, 93, 135, 21, 161, 150,
 41, 16, 123, 154, 199, 243, 145, 120, 111, 157, 158,
 178, 177, 50, 117, 25, 61, 255, 53, 138, 126, 109,
 84, 198, 128, 195, 189, 13, 87, 223, 245, 36, 169,
 62, 168, 67, 201, 215, 121, 214, 246, 124, 34, 185,
 3, 224, 15, 236, 222, 122, 148, 176, 188, 220, 232,
 40, 80, 78, 51, 10, 74, 167, 151, 96, 115, 30,
 0, 98, 68, 26, 184, 56, 130, 100, 159, 38, 65,
 173, 69, 70, 146, 39, 94, 85, 47, 140, 163, 165,
 125, 105, 213, 149, 59, 7, 88, 179, 64, 134, 172,
 29, 247, 48, 55, 107, 228, 136, 217, 231, 137, 225,
 27, 131, 73, 76, 63, 248, 254, 141, 83, 170, 144,
 202, 216, 133, 97, 32, 113, 103, 164, 45, 43, 9,
 91, 203, 155, 37, 208, 190, 229, 108, 82, 89, 166,
 116, 210, 230, 244, 180, 192, 209, 102, 175, 194, 57,
 75, 99, 182).

4.2. Linear Transformation

The linear transformation is denoted by $l: (V_8)^{16} \rightarrow V_8$, and defined as:

$$l(a_{15}, \dots, a_0) = \text{nabla}(148 \cdot \delta(a_{15}) + 32 \cdot \delta(a_{15}) + 133 \cdot \delta(a_{13}) + 16 \cdot \delta(a_{12}) + 194 \cdot \delta(a_{11}) + 192 \cdot \delta(a_{10}) + 1 \cdot \delta(a_9) + 251 \cdot \delta(a_8) + 1 \cdot \delta(a_7) + 192 \cdot \delta(a_6) + 194 \cdot \delta(a_5) + 16 \cdot \delta(a_4) + 133 \cdot \delta(a_3) + 32 \cdot \delta(a_2) + 148 \cdot \delta(a_1) + 1 \cdot \delta(a_0)),$$

for all a_i belonging to V_8 , $i = 0, 1, \dots, 15$, where the addition and multiplication operations are in the field Q , and constants are elements of the field as defined above.

4.3. Transformations

The following transformations are applicable for encryption and decryption algorithms:

$X[x]: V_{128} \rightarrow V_{128}$ $X[k](a) = x(xor)a$, where k, a belong to V_{128} ,

$S:V_{128} \rightarrow V_{128}$ $S(a)=(a_{15}||\dots||a_0)=\pi(a_{15})||\dots||\pi(a_0)$, where $a_{15}||\dots||a_0$ belongs to V_{128} , a_i belongs to V_8 , $i=0,1,\dots,15$,

$S^{(-1)}:V_{128} \rightarrow V_{128}$ the inverse transformation of S , which may be calculated, for example, as follows:
 $S^{(-1)}(a_{15}||\dots||a_0)=\pi^{(-1)}(a_{15})||\dots||\pi^{(-1)}(a_0)$, where $a_{15}||\dots||a_0$ belongs to V_{128} , a_i belongs to V_8 , $i=0,1,\dots,15$, $\pi^{(-1)}$ is the inverse of π .

$R:V_{128} \rightarrow V_{128}$ $R(a_{15}||\dots||a_0)=l(a_{15},\dots,a_0)||a_{15}||\dots||a_1$, where $a_{15}||\dots||a_0$ belongs to V_{128} , a_i belongs to V_8 , $i=0,1,\dots,15$,

$L:V_{128} \rightarrow V_{128}$ $L(a)=R^{(16)}(a)$, where a belongs to V_{128} ,

$R^{(-1)}:V_{128} \rightarrow V_{128}$ the inverse transformation of R , which may be calculated, for example, as follows: $R^{(-1)}(a_{15}||\dots||a_0)=a_{14}||a_{13}||\dots||a_0||l(a_{14},a_{13},\dots,a_0,a_{15})$, where $a_{15}||\dots||a_0$ belongs to V_{128} , a_i belongs to V_8 , $i=0,1,\dots,15$, $\pi^{(-1)}$ is the inverse of π ,

$L^{(-1)}:V_{128} \rightarrow V_{128}$ $L^{(-1)}(a)=(R^{(-1)})^{(16)}(a)$, where a belongs to V_{128} ,

$F[k]:V_{128}[*]V_{128} \rightarrow V_{128}[*]V_{128}$
 $F[k](a_1,a_0)=(LSX[k](a_1)(\text{xor})a_0,a_1)$, where k , a_0 , a_1 belong to V_{128} .

4.4. Key schedule

Key schedule uses round constants C_i belonging to V_{128} , $i=1, 2, \dots, 32$, defined as

$$C_i=L(\text{Vec}_{128}(i)), i=1,2,\dots,32.$$

Round keys K_i , $i=1, 2, \dots, 10$ are derived from key $K=k_{255}||\dots||k_0$ belonging to V_{256} , k_i belongs to V_1 , $i=0, 1, \dots, 255$, as follows:

$$\begin{aligned} K_1 &= k_{255}||\dots||k_{128}; \\ K_2 &= k_{127}||\dots||k_0; \\ (K_{(2i+1)}, K_{(2i+2)}) &= F[C_{(8(i-1)+8)}] \dots F[C_{(8(i-1)+1)}] \\ & \quad (K_{(2i-1)}, K_{(2i)}), i=1,2,3,4. \end{aligned}$$

4.5. Basic encryption algorithm

79d26221b87b584cd42fbc4ffea5de9a,
L(79d26221b87b584cd42fbc4ffea5de9a) =
0e93691a0cfc60408b7b68f66b513c13,
L(0e93691a0cfc60408b7b68f66b513c13) =
e6a8094fee0aa204fd97bcb0b44b8580.

5.4. Key schedule

In this test example, the key is equal to:

```
K =
8899aabbccddeeff0011223344556677fedcba98765432100123456789abcdef.

K_1 = 8899aabbccddeeff0011223344556677,
K_2 = fedcba98765432100123456789abcdef.

C_1 = 6ea276726c487ab85d27bd10dd849401,
X[C_1](K_1) = e63bdcc9a09594475d369f2399d1f276,
SX[C_1](K_1) = 0998ca37a7947aabb78f4a5ae81b748a,
LSX[C_1](K_1) = 3d0940999db75d6a9257071d5e6144a6,
F[C_1](K_1, K_2) = (c3d5fa01ebe36f7a9374427ad7ca8949,
8899aabbccddeeff0011223344556677).

C_2 = dc87ece4d890f4b3ba4eb92079cbeb02,
F[C_2]F[C_1](K_1, K_2) = (37777748e56453377d5e262d90903f87,
c3d5fa01ebe36f7a9374427ad7ca8949).

C_3 = b2259a96b4d88e0be7690430a44f7f03,
F[C_3]...F[C_1](K_1, K_2) = (f9eae5f29b2815e31f11ac5d9c29fb01,
37777748e56453377d5e262d90903f87).

C_4 = 7bcd1b0b73e32ba5b79cb140f2551504,
F[C_4]...F[C_1](K_1, K_2) = (e980089683d00d4be37dd3434699b98f,
f9eae5f29b2815e31f11ac5d9c29fb01).

C_5 = 156f6d791fab511deabb0c502fd18105,
F[C_5]...F[C_1](K_1, K_2) = (b7bd70acea4460714f4ebe13835cf004,
e980089683d00d4be37dd3434699b98f).

C_6 = a74af7efab73df160dd208608b9efe06,
F[C_6]...F[C_1](K_1, K_2) = (1a46ea1cf6ccd236467287df93fdf974,
b7bd70acea4460714f4ebe13835cf004).

C_7 = c9e8819dc73ba5ae50f5b570561a6a07,
F[C_7]...F[C_1](K_1, K_2) = (3d4553d8e9cfec6815ebadc40a9ffd04,
1a46ea1cf6ccd236467287df93fdf974).

C_8 = f6593616e6055689adfb18027aa2a08,
(K_3, K_4) = F[C_8]...F[C_1](K_1, K_2) =
(db31485315694343228d6aef8cc78c44, 3d4553d8e9cfec6815ebadc40a9ffd04).
```

The round keys K_i , $i = 1, 2, \dots, 10$, take the following values:

```
K_1 = 8899aabbccddeeff0011223344556677,
K_2 = fedcba98765432100123456789abcdef,
K_3 = db31485315694343228d6aef8cc78c44,
K_4 = 3d4553d8e9cfec6815ebadc40a9ffd04,
```

K_5 = 57646468c44a5e28d3e59246f429f1ac,
K_6 = bd079435165c6432b532e82834da581b,
K_7 = 51e640757e8745de705727265a0098b1,
K_8 = 5a7925017b9fdd3ed72a91a22286f984,
K_9 = bb44e25378c73123a5f32f73cdb6e517,
K_10 = 72e9dd7416bcf45b755dbaa88e4a4043.

5.5. Test encryption

In this test example, encryption is performed on the round keys specified in clause 5.4. Let the plaintext be

$$a = 1122334455667700\text{ffeeddccbbaa}9988,$$

then

$$\begin{aligned} X[K_1](a) &= 99\text{bb}99\text{ff}99\text{bb}99\text{ffffff}\text{ffffff}\text{ffffff}\text{ffffff}, \\ SX[K_1](a) &= \text{e}87\text{de}8\text{b}6\text{e}87\text{de}8\text{b}6\text{b}6\text{b}6\text{b}6\text{b}6\text{b}6\text{b}6\text{b}6, \\ LSX[K_1](a) &= \text{e}297\text{b}686\text{e}355\text{b}0\text{a}1\text{c}f4\text{a}2\text{f}9249140830, \\ LSX[K_2]LSX[K_1](a) &= 285\text{e}497\text{a}0862\text{d}596\text{b}36\text{f}4258\text{a}1\text{c}69072, \\ LSX[K_3]\dots LSX[K_1](a) &= 0187\text{a}3\text{a}429\text{b}567841\text{ad}50\text{d}29207\text{cc}34\text{e}, \\ LSX[K_4]\dots LSX[K_1](a) &= \text{ec}9\text{b}d\text{ba}057\text{d}4\text{f}4\text{d}77\text{c}5\text{d}70619\text{dcad}206, \\ LSX[K_5]\dots LSX[K_1](a) &= 1357\text{f}d11\text{de}9257290\text{c}2\text{a}1473\text{eb}6\text{b}c\text{de}1, \\ LSX[K_6]\dots LSX[K_1](a) &= 28\text{ae}31\text{e}7\text{d}4\text{c}2354261027\text{ef}0\text{b}32897\text{df}, \\ LSX[K_7]\dots LSX[K_1](a) &= 07\text{e}223\text{d}56002\text{c}013\text{d}3\text{f}5\text{e}6\text{f}714\text{b}86\text{d}2\text{d}, \\ LSX[K_8]\dots LSX[K_1](a) &= \text{cd}8\text{ef}6\text{cd}97\text{e}0\text{e}092\text{a}8\text{e}4\text{cca}61\text{b}38\text{bf}65, \\ LSX[K_9]\dots LSX[K_1](a) &= 0\text{d}8\text{e}40\text{e}4\text{a}800\text{d}06\text{b}2\text{f}1\text{b}37\text{ea}379\text{ead}8\text{e}. \end{aligned}$$

Then the ciphertext is

$$b = X[K_{10}]LSX[K_9]\dots LSX[K_1](a) = 7\text{f}679\text{d}90\text{bebc}24305\text{a}468\text{d}42\text{b}9\text{d}4\text{edcd}.$$

5.6. Test decryption

In this test example, decryption is performed on the round keys specified in clause 5.4. Let the ciphertext be

$$b = 7\text{f}679\text{d}90\text{bebc}24305\text{a}468\text{d}42\text{b}9\text{d}4\text{edcd},$$

then

$$\begin{aligned} X[K_{10}](b) &= 0\text{d}8\text{e}40\text{e}4\text{a}800\text{d}06\text{b}2\text{f}1\text{b}37\text{ea}379\text{ead}8\text{e}, \\ L^{(-1)}X[K_{10}](b) &= 8\text{a}6\text{b}930\text{a}52211\text{b}45\text{c}5\text{baa}43\text{ff}8\text{b}91319, \\ S^{(-1)}L^{(-1)}X[K_{10}](b) &= 76\text{ca}149\text{eef}27\text{d}1\text{b}10\text{d}17\text{e}3\text{d}5\text{d}68\text{e}5\text{a}72, \\ S^{(-1)}L^{(-1)}X[K_9]S^{(-1)}L^{(-1)}X[K_{10}](b) &= 5\text{d}9\text{b}06\text{d}41\text{b}9\text{d}1\text{d}2\text{d}04\text{df}7755363\text{e}94\text{a}9, \\ S^{(-1)}L^{(-1)}X[K_8]\dots S^{(-1)}L^{(-1)}X[K_{10}](b) &= 79487192\text{aa}45709\text{c}115559\text{d}6\text{e}9280\text{f}6\text{e}, \\ S^{(-1)}L^{(-1)}X[K_7]\dots S^{(-1)}L^{(-1)}X[K_{10}](b) &= \text{ae}506924\text{c}8\text{ce}331\text{bb}918\text{fc}5\text{bdfb}195\text{fa}, \\ S^{(-1)}L^{(-1)}X[K_6]\dots S^{(-1)}L^{(-1)}X[K_{10}](b) &= \text{bbffbf}c8939\text{eaaffaf}b8\text{e}22769\text{e}323\text{aa}, \\ S^{(-1)}L^{(-1)}X[K_5]\dots S^{(-1)}L^{(-1)}X[K_{10}](b) &= 3\text{cc}2\text{f}07\text{cc}07\text{a}8\text{bec}0\text{f}3\text{ea}0\text{ed}2\text{ae}33\text{e}4\text{a}, \\ S^{(-1)}L^{(-1)}X[K_4]\dots S^{(-1)}L^{(-1)}X[K_{10}](b) &= \text{f}36\text{f}01291\text{d}0\text{b}96\text{d}591\text{e}228\text{b}72\text{d}011\text{c}36, \\ S^{(-1)}L^{(-1)}X[K_3]\dots S^{(-1)}L^{(-1)}X[K_{10}](b) &= \end{aligned}$$

1c4b0c1e950182b1ce696af5c0bfc5df,
 $S^{(-1)}L^{(-1)}X[K_2] \dots S^{(-1)}L^{(-1)}X[K_{10}](b) =$
99bb99ff99bb99ffffffffffffffffffff.

Then the plaintext is

$$a = X[K_1]S^{(-1)}L^{(-1)}X[K_2]\dots S^{(-1)}L^{(-1)}X[K_{10}](b) =$$

1122334455667700ffeeddccbbaa9988.

6. Security Considerations

This entire document is about security considerations.

7. References

7.1. Normative References

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